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Легковесное профилирование разделяемых библиотек в Linux

для встраиваемых систем

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Motivation

Popular Linux Profilers (gprof, gcov, GPT, Valgrind) have following **problems**:

- Need to recompile with special options (gprof, gcov)
- Need to **relink** program with 3rd-party libraries (GPT)
- Need to use special environment for profiling (Valgrind)
- Need to use superuser rights

Project goals

We need a tool for system-wide ELF executables performance analysis.

This tool should **allow user** next things **do easily**:

- Profile function calls without recompilation and relinking with 3rd-party libraries
- Profile only given set of C/C++ functions from shared libraries in Linux
- Profile both dynamically linked and dynamically loaded functions
- Profile without creating of special environment
- Get information about number and total duration of function calls
- Perform profiling on x86/x64 platforms

"Non-invasive" profiling

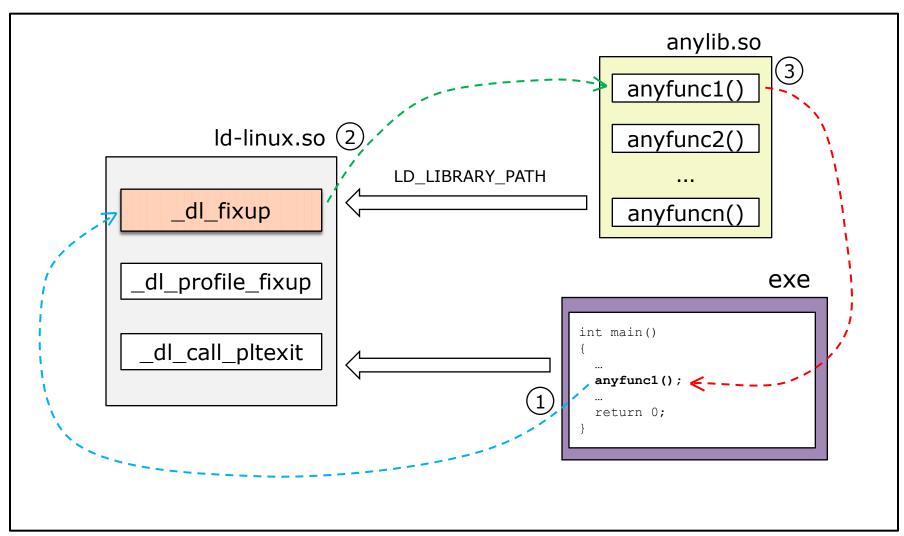
Main ideas:

- Profiler can not be implemented in the program code
- Profiling should be performed at well-defined points of function calls
- Profiling process should not corrupt the algorithm of profiled application
- Profiling process should use minimum amount of system resources
- Results of profiling should be as accurate as possible

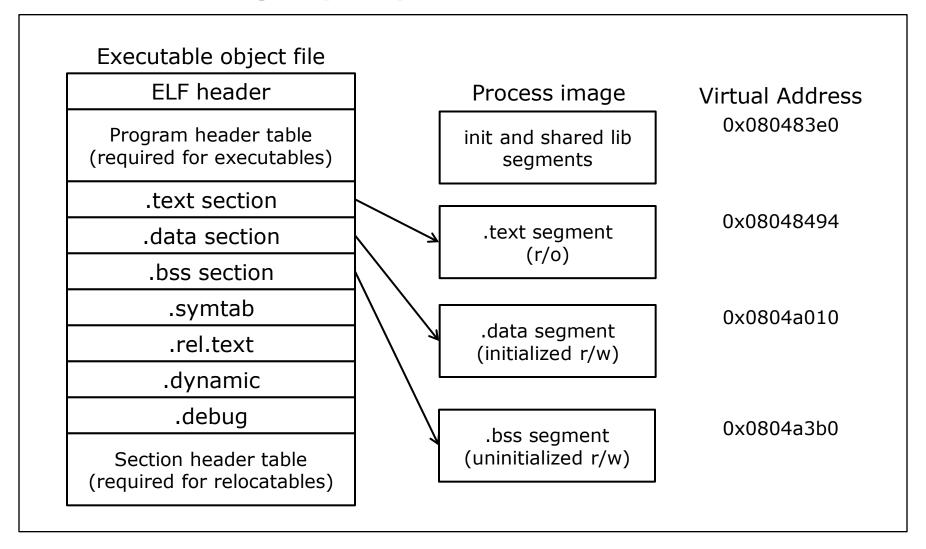
Ways to implement

- Infiltration into the symbol relocation process
- Modification of Linux dynamic linker (ld-linux.so)
- Modification of dynamic loading library (libdl.so)

Dynamic Linking



ELF Parsing by Dynamic Linker



1. Resolving the Dependencies

- When linking a dynamic executable, one or more shared objects are explicitly referenced. These objects are recorded as dependencies within the dynamic executable.
- The runtime linker uses this dependency information to locate, and load, the associated objects.
- Once all the dynamic executable's dependencies are loaded, each dependency is inspected, in the order the dependency is loaded, to locate any additional dependencies.

1. Resolving the Dependencies

- The Linux runtime linker looks in two default locations for dependencies /lib and /usr/lib.
- The dependencies of a dynamic executable or shared object can be displayed using **Idd**. For example, the file /usr/bin/cat has the following dependencies:

\$ Idd /usr/bin/cat

```
libc.so.1 => /lib/libc.so.1 libm.so.2 => /lib/libm.so.2
```

• The dependencies recorded in an object can be inspected using **dump**. Use this command to display the file's .dynamic section, and look for entries that have a NEEDED tag.

```
$ dump -Lvp prog
```

```
prog:
```

[INDEX] Tag Value

[1] NEEDED libfoo.so.1

[2] NEEDED libc.so.1

[3] RUNPATH /home/me/lib:/home/you/lib

.

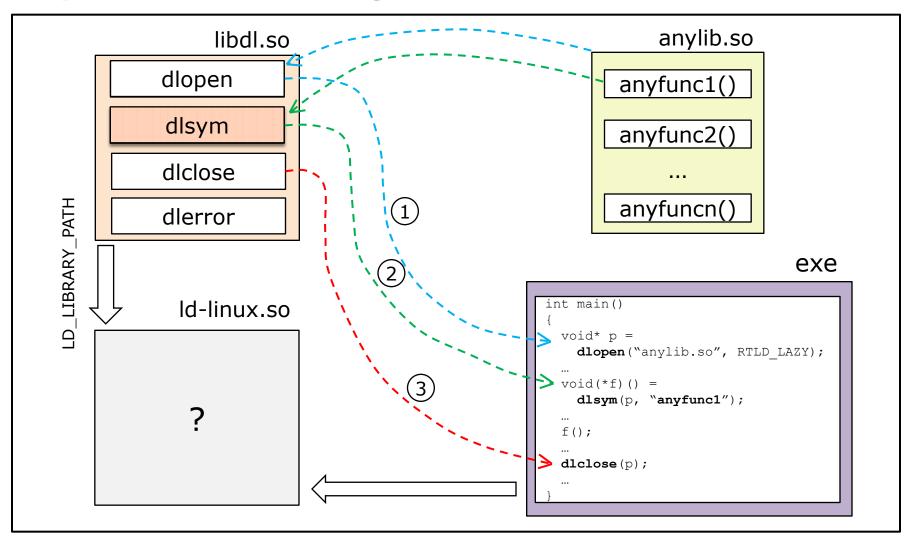
Symbol Table Structure

```
typedef struct {
    Elf32_Word st_name;
    Elf32_Addr st_value;
    Elf32_Word st_size;
    unsigned char st_info;
    unsigned char st_other;
    Elf32_Half st_shndx;
} Elf32_Sym;
```

Parsing other sections of ELF

- For dynamic linking, the Dynamic linker primarily uses two processor-specific tables:
 - Global Offset Table (GOT)
 - Procedure Linkage Table (PLT)
- Dynamic linkers support PIC Code through the GOT in each shared library
- The GOT contains absolute addresses to all of the static data referenced in the program.

Dynamic Loading



Profiler components

- Shared library libelfperf.so
 - Call redirection and function wrapping mechanisms
 - Collecting of calls statistics
 - Memory management
- Modified dynamic linker (ld-linux.so)
 - Uses libelfperf.so for profiling of dynamically linked functions
 - Displays the results of profiling
- Modified dynamic loading library (libdl.so)
 - Uses libelfperf.so for profiling of dynamically loaded functions

Call redirection mechanism

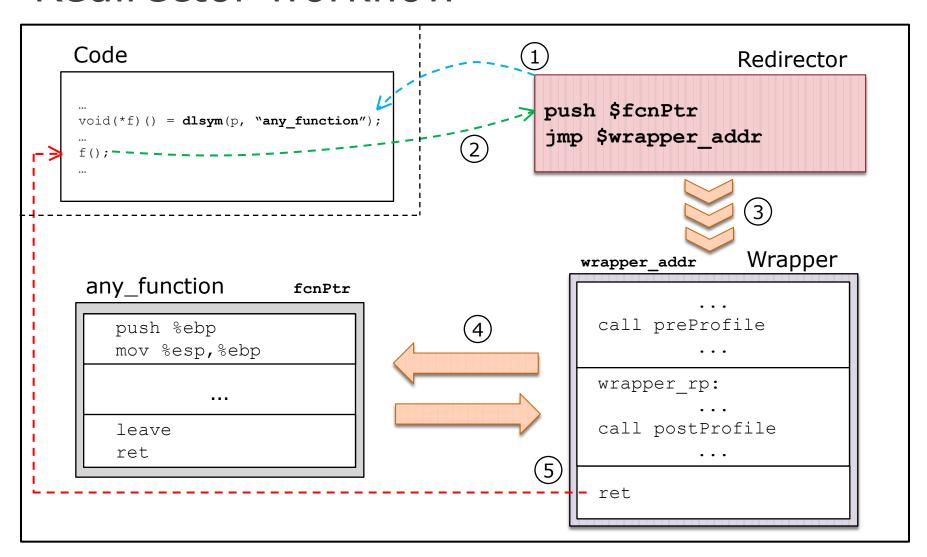
Calls redirection mechanism (**Redirector**) is a set of machine codes for the next assembly instructions:

```
push $fcnPtr
jmp $wrapper_addr
```

All they do is:

- Save address of profiled function in program stack
- Jump to wrapper-function

Redirector workflow



Redirector details

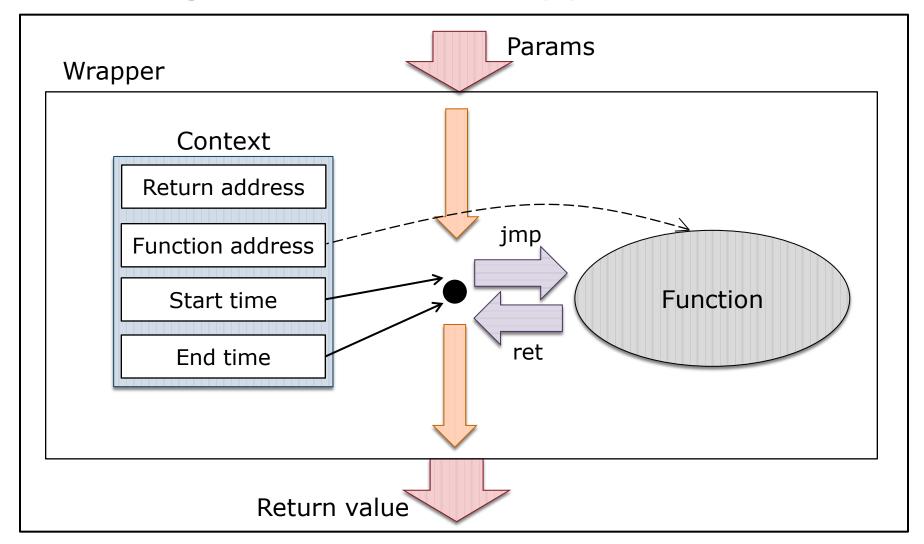
- Each redirector is created individually for each profiled function
- Redirectors are placed into data segment of process virtual memory
- The operating system allows to mark these memory areas as executable

Wrapping mechanism

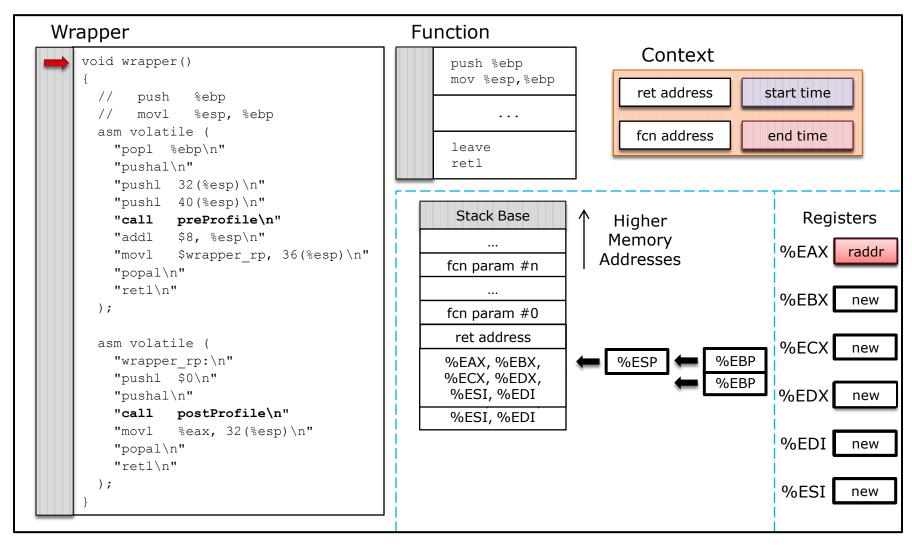
Function Wrapping mechanism (or **Wrapper**) is a function that does next things:

- Takes control from redirector
- Performs pre-profile operations
- Performs replacement of return address
- Performs jump into profiled function
- Again takes control after the work of profiled function
- Performs post-profile operations
- Returns to caller

Working scheme of Wrapper



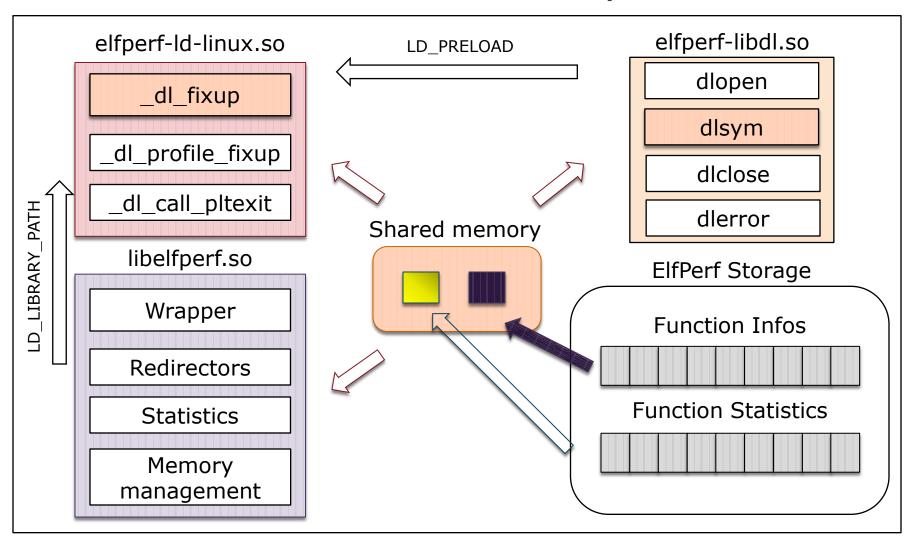
Implementation details (x86)



Wrapper details

- Wrapper doesn't corrupt stack content
- Wrapper exists in a single copy for all functions in each profiler implementation (x86 or x64)
- Saving/Restoring of registers' state allows to escape of uncontrollable changes in the program state
- Allows to profile wide set of C/C++ functions

Interaction of ElfPerf's components



Conclusion

Now we have:

- *Light* profiler based on *patched* Id-linux.so and libdl.so
- Support of profiling for C/C++ functions from shared libraries (including libs compiled with -fomit-frame-pointer flag)
- Collecting of information about number and total duration of function calls
- Support of both x86 and x64 platforms

Links

Project resources:

- https://github.com/OSLL/elfperf
- http://dev.osll.ru/projects/epat/wiki/

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